





FIRST FRAMEWORK ON SHaRK OS



Giuseppe Lipari, Michael Trimarchi RETIS Lab Scuola Superiore Sant'Anna

Scuola
The Superiore S.
University of Anna of Pisa
Cantabria



Summary

Software Framework Status of implementation in Shark Examples of usage of the API

Creating a contract for

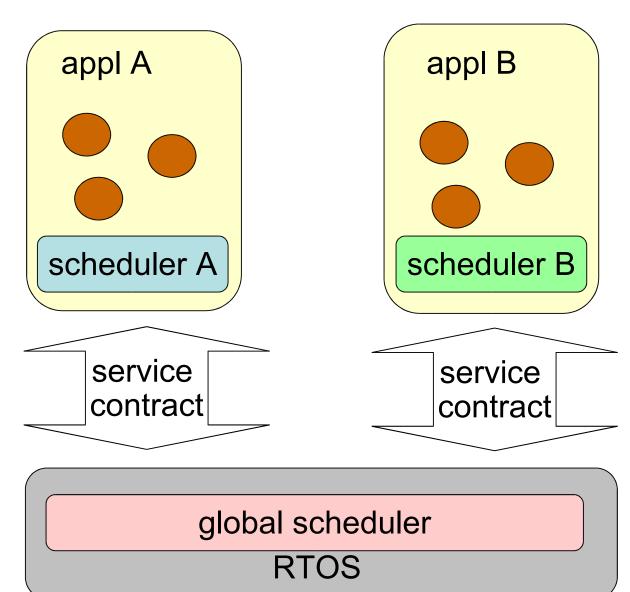
- periodic, sporadic hard real-time tasks
- soft real-time task
- imprecise computation
- applications (set of tasks)



Software Framework

- Application = set of tasks/threads (+ scheduler)
 - Synonyms: Component, Subsystem
 - It can be reduced to one single task
- Hierarchical scheduling structure
 - System = set of applications
 - Each application may have its own local scheduler
- Service Contract
 - Each application specifies its requirements by requiring a service contract







Server based scheduling

Server-based scheduling

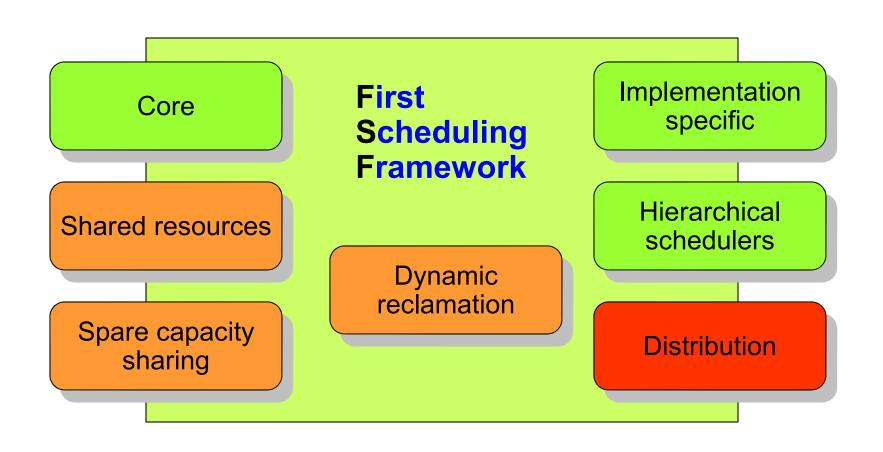
- Each application is assigned one or more servers
- Each server has a budget and a period
- Provides temporal isolation
- Provides independent analysis

Server algorithm

- No specific global scheduling strategy
- No specific server algorithm
- Systems can be based on
 - Fixed Priority and Sporadic Server
 - EDF and Constant Bandwidth Server
 - Table Driven and Slot Shifting



Status of implementation in Shark





Service contract parameters

Basic attributes

Budget_min, max period_min, max workload_type

Timing Attributes

d_equals_t, dline, budget_overrun dline_overrun

Hierarchical sched.

type

Shared resources

list of prot. operations resource id operation id

Implementation spec.

server preempt. lev. resource preempt. lev.

Spare cap. sharing

granularity, util. set quality & importance



Core service contract



Budget_min, max period_min, max workload_type

Timing Attributes

d_equals_t, dline, budget_overrun dline_overrun

Hierarchical sched.

type

Shared resources

list of prot. operations resource id operation id

Implementation spec.

server preempt. lev. resource preempt. lev.

Capacity sharing

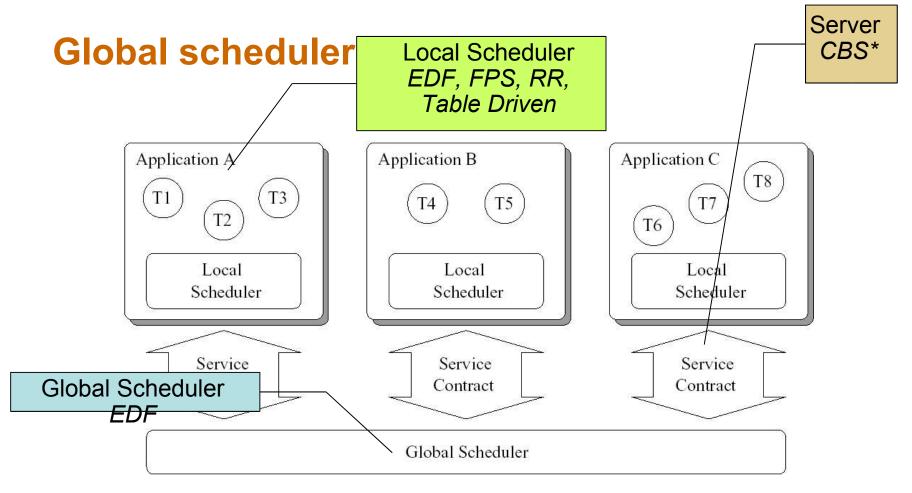
granularity, util. set quality & importance

Basic server mechanism: Constant Bandwidth Server (CBS)

Basic scheduling mechanism: EDF

Negotiation mechanism: polynomial schedulability test



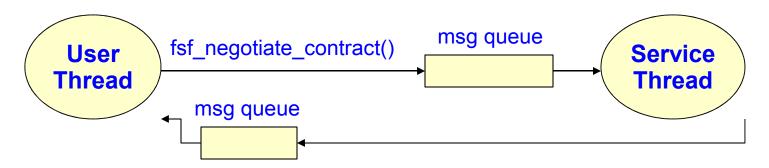


No implementation-specific data is needed

- no preemption level for servers (automatically assigned by EDF)
- preemption level for resources is not needed either



Contract Negotiation

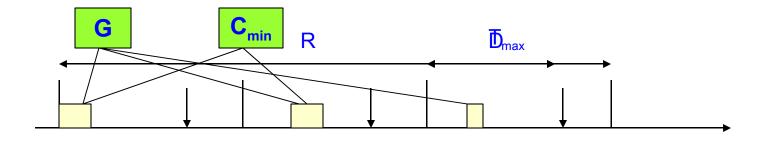


Client/server structure:

the service thread is assigned a contract response time:

- Service thread contract params: (C_{min}, T_{max}, D)
- Computation time for negotiation: G

$$R = \left\lfloor \frac{G}{C_{\min}} \right\rfloor T_{\max} + D$$





Synchronization

Synchronization between servers

 When two threads in two different servers use the same mutex

Synchronization mechanism

- we are using a mechanism called BWI (Bandwidth inheritance)
- Similar to Priority Inheritance

Shared resources

list of prot. operations resource id operation id

Implementation spec.

server preempt. lev. resource preempt. lev.

Capacity sharing

granularity, util. set quality & importance



Hierarchical scheduling

Basic attributes

Budget_min, max period_min, max workload_type

Timing Attributes

d_equals_t, dline, budget_overrun dline_overrun

Hierarchical sched.

type

Current support for

- Fixed Priority (POSIX std)
- Round Robin (POSIX std)
- EDF
- Table Driven (with deadline transformation)

Easy to introduce new schedulers

- thanks to Shark modularity
- not part of the API



Capacity sharing

The contract is flexible

- possibility of re-negotiation
- possibility of obtaining more than the minimum

In Shark

- Use elastic task (Buttazzo et al.) to assign spare capacity
- Among those with equal importance, the quality parameter is used as the elastic constant

Shared resources

list of prot. operations resource id operation id

Implementation spec.

server preempt. lev. resource preempt. lev.

Capacity sharing

granularity, util. set quality & importance

Not completed:

- this feature is to be used only when D=T
- planning extension to general model in next phase



Dynamic reclaimation

If some thread execute less than expected, the spare capacity is dynamically reassigned

Current implementation

- the GRUB (greedy reclaimation of unused bandwidth) has been implemented in Shark
- it is not possible to specify which thread gets the extra capacity
- no parameter in the interface

To be done

- the algorithm is valid if D=T
- to be extended to the general model



Examples of usage of the API

Example 1: Initialize a contract for single thread Define a contract fsf contract parameters t contract; fsf server id t server; pthread t j; Initizalize the contract fsf initialize contract(&contract); fsf set contract basic parameters (&contract, &cmin, &tmax, &cmax, &tmin, workload); fsf_set_contract_timing_requirements(&contract,FALSE,&deadline,0, no sigval, 0, no sigval); if (!fsf negotiate contract(&contract,&server)) { // ERROR else fsf_create_thread(servex,&j,NULL,task,NULL,NULL); Create the thread Negotiate and bind it to the server



Example: typical thread structure

Periodic thread

```
void task body(void *arg)
  struct timespec acttime;
  struct timespec budget;
  struct timespec
                     period;
                       deadline missed;
  bool
                       budget overrun;
  bool
  int
                       uperiod;
  [...]
  sys gettime (&acttime);
  while (1)
       ADDUSEC2TIMESPEC (uperiod, &acttime);
       fsf schedule next timed job (&acttime, &budget, &period,
                &budget overrun, &deadline missed);
        /* Bodv */
```



Example: typical thread structure (2)

Aperiodic thread

```
void task body(void *arg) {
                                                     Synchronization
                                                     object
  fsf synch object handle t synch handle;
  struct timespec
                         budget;
  struct timespec
                         period;
  bool
                          deadline missed;
  bool
                          budget overrun;
  [...]
                                                    Wait for
  while(1) {
                                                    next synch.
    fsf_schedule_next_event_triggered job &synch handle, &budget,
             &period, &budget overrun, &deadline missed);
     * Body */
```



Example: hard real-time periodic threads

What is needed

Core service (+ Shared resource synchronization)

Contract Parameters

 $C_{min} = C_{max} = WCET$ of the thread

 $T_{min} = T_{max} = thread's period$

D = thread's deadline

workload = bounded

budget overrun exception handling

Advantages

The thread is protected from the other non-RT and soft RT threads in the system (temporal isolation)

if dynamic reclaimation, the spare capacity of this thread can be given to others



Example: soft real-time periodic threads

What is needed

Core + (capacity sharing) + (dynamic recl.) + (shared res. synch.)

Contract Parameters

 $C_{min} - C_{max} = variation of the execution time$

 $T_{min} = T_{max} = thread's period$

D = thread's deadline

workload = indeterminate

Advantages

Does not impact on other threads (temporal isolation)

minimum service is guaranteed

Takes advantage of capacity sharing and dynamic reclamation (to minimize deadline misses)

can re-negotiate if it needs more



Example: imprecise computation

Thread consists of a mandatory part and N optional parts

- WCET of mandatory part = M
- WCET of i-th optional part = O_i

What is needed

Core + (capacity sharing) + (dynamic recl.) + (shared res. synch.)

Contract Parameters

$$C_{\min} = M$$

$$C_{max} = M + O_1 + ... + O_N$$

$$T_{min} = T_{max} =$$
thread's period

D = thread's deadline

workload = bounded



Example: thread structure of an imprecise computation thread

```
void task body(void *arg) {
  pthread t my pid = (pthread t) (arg);
  int i;
  sys gettime (&acttime);
                                                               Get remaining
  fsf get server (&server, my pid);
                                                               capacity
  while(1) {
    /* Mandatory Body */
    for (i=0; i< N; i++) {
        fsf get available capacity (server, &capa\dot{c}ity);
        if \overline{\text{capacity}} > O[\overline{1}]
         /* Optional Code */
                                                          If enough capacity
        else break;
                                                          execute optional
    ADDUSEC2TIMESPEC (uperiod, &acttime);
                                                          part
    fsf schedule next timed job(&acttime, &budget, &period,
                  &budget missed, &deadline missed);
```



Example: imprecise computation

Advantage:

mandatory part is guaranteed

If capacity sharing and dynamic reclamation services are available, some optional part may be completed as well

More reclaiming 🗗 🕭 more optional parts



Example: creating a contract for an application

What is needed

Core + Hierarchical

Contract Parameters

workload = indeterminate

Scheduler type = EDF or FPS or RR or TD

Advantages

re-using an existing code base without re-designing and reimplementing it

need minimal modifications to the original code



Example: contract for hierarchical

```
fsf contract parameters t contract;
fsf server id t server;
                                                                  Set scheduler
pthread t j, k;
HARD TASK MODEL ht;
fsf initialize contract(&contract);
fsf set contract basic parameters (&contract, &cmin, &tmax,
                           &cmax, &tmin, workload);
fsf set contract timing requirements (&contract, FALSE, &deadline, 0,
                           no siqual, 0, no siqual);
fsf set local scheduler parameter (&contract, FSF SCHEDULER EDF);
fsf negotiate contract(&contract,&server);
/* S.Ha.R.K. hard task parameters */
hard task default model(ht);
hard task def mit (ht, TIMESPEC2USEC (&deadline));
hard task def wcet(ht, TIMESPEC2USEC(&wcet));
                                                            Create a local thread
/* Create EDF task */
fsf create local thread(server,&j,NULL,task,NULL,&ht);
/* Create EDF task */
fsf create local thread(server, &k, NULL, task, NULL, &ht);
```

